How to Use Mean-Field Control for Restless Bandits and Weakly Coupled MDPs

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joint work with Bruno Gaujal, Kimang Khun, Chen Yan

Inria

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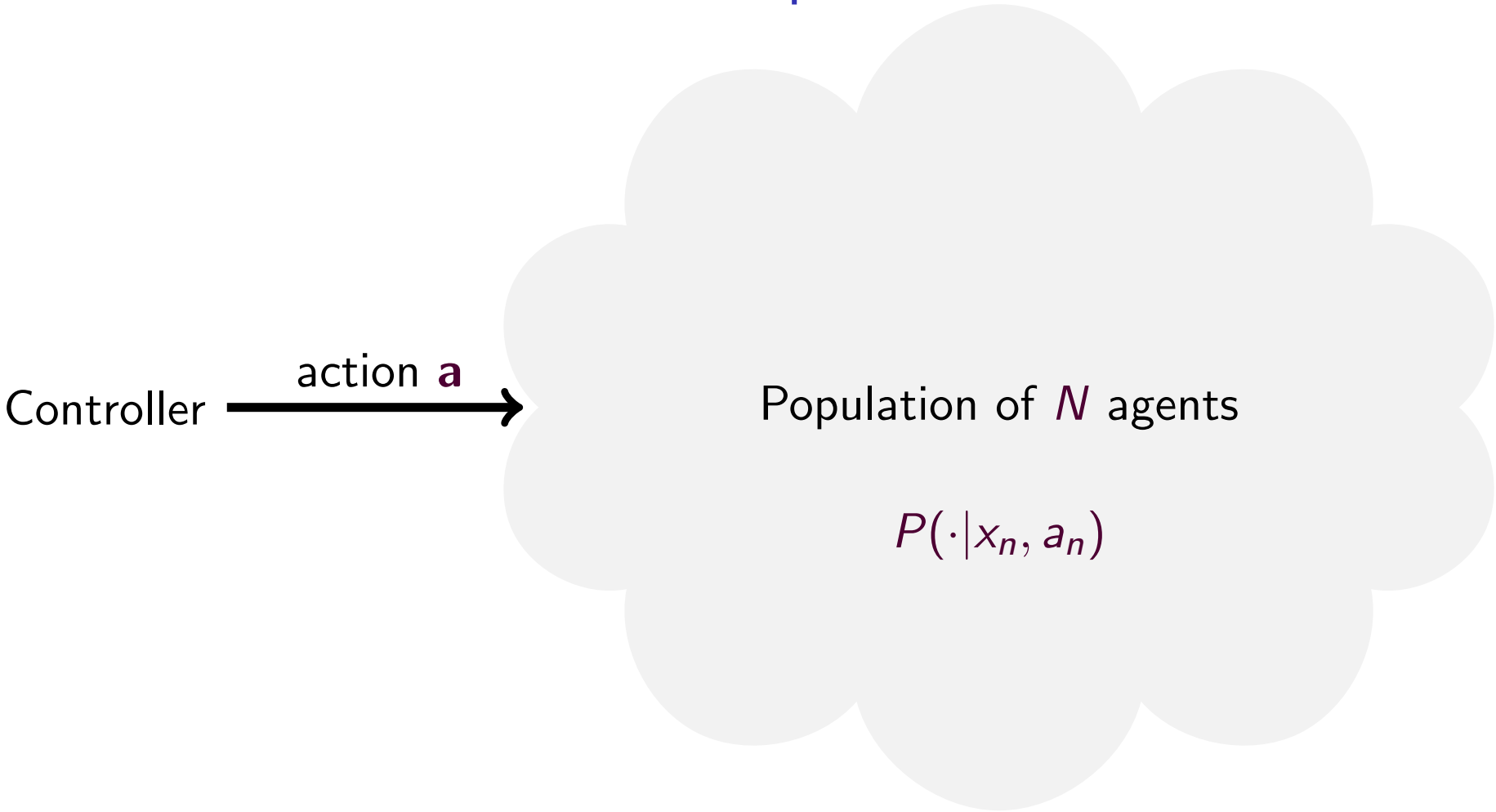
Centralized mean field control problem

Controller $\xrightarrow{\text{action } \mathbf{a}}$ Population of N agents

$P(\cdot|x_n, a_n)$

Centralized mean field control problem



The computational difficulty increases with N but $N = \infty$ is easy.

- How to use the $N = +\infty$ solution for finite N?
- How efficient is this? (i.e., how fast does it become optimal?)

This talk will focus on *Markovian bandits*

N statistically identical arms (=agents)

- Discrete time, finite state space.
- $P(\cdot|s_n, a_n)$ and $r(s_n, a_n)$.

Maximize expected reward

$$\frac{1}{T}\sum_{t=1}^{T}\sum_{n=1}^{N} r(s_n(t), a_n(t)).$$

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Hard constraint: $\forall t : \sum_{n=1}^{N} a_n(t) \leq C.$

- If $a_n(t) \in \{0,1\}$: Markovian bandit (this talk)
- If $a_n(t) \in \{0,1\}^d$: Weakly coupled MDP.

Example 1: Applicant screening problem

N applicants, T rounds of interview.
Each round: you can interview up to αN candidates.
Goal: maximize the expected quality of selected candidates.

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Each candidate has an (unknown) quality p_n.

- Result of an interview: Bernoulli(p_n)

Goal: find the βN highest p_n.

Example 2: What to work on?

Job Scheduling

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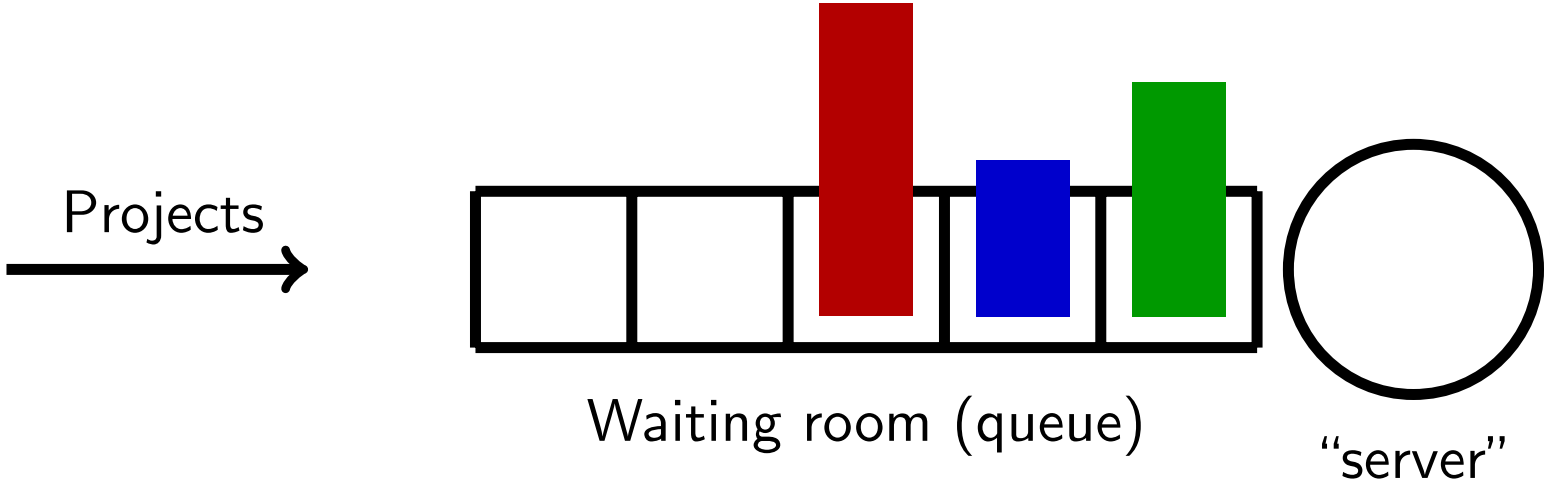

- Examples: research projects, tasks allocations, electric vehicle charging, wireless scheduling,...
- Heuristics: SRPT, EDF,...

Main questions and outline

These problems are restless bandit problems (PSPACE-hard)

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- Close-to-optimal (if possible)

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- Close-to-optimal (if possible)

Questions

- When are priority rules asymptotically optimal?
- How to compute a good priority rule?
- How fast do they become optimal?

Outline

1. The $N = +\infty$ problem
2. Infinite-horizon and index policies
3. Asymptotic optimality and index computation
4. Finite-horizon restless bandits
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The $N = \infty$ problem

Original model: $P(\cdot|s_n, a_n)$ and $r(s_n, a_n)$. For all t, $\sum_{n=1}^{N} a_n(t) \leq \alpha N$.

The $N = \infty$ problem

Relaxed model: $P(\cdot|s_n, a_n)$ and $r(s_n, a_n)$. For all t, $\mathbb{E}\left[\sum_{n=1}^{N} a_n(t)\right] \leq \alpha N$.

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- $x_s = \mathbf{P}[s_n = s]$ and $y_{s,a} = \mathbf{P}[s_n = s, a_n = a]$.

$$\max_{x \geq 0, y \geq 0} \sum_{s,a} r_{s,a} y_{s,a}$$

$$\text{s.t.} \quad x_{s'} = \sum_{s} y_{s,a} P(s'|s,a)$$

$$x_s = \sum_{a} y_{s,a}$$

$$\sum_{s} x_s = 1.$$

$$\sum_{s} y_{s,1} = \alpha \qquad \text{relaxed budget contraint}$$

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- $x_s(t) = \mathbf{P}\left[s_n(t) = s\right]$ and $y_{s,a}(t) = \mathbf{P}\left[s_n(t) = s, a_n(t) = a\right]$.

$$\max_{x \geq 0, y \geq 0} \sum_{t=1}^{T} \sum_{s,a} r_{s,a} y_{s,a}(t)$$

$$\text{s.t.} \quad x_{s'}(t+1) = \sum_{s} y_{s,a}(t) P(s'|s,a)$$

$$x_s(t) = \sum_{a} y_{s,a}(t)$$

$$\sum_{s} x_s = x_s(0).$$

$$\sum_{s} y_{s,1}(t) = \alpha(t) \qquad \text{relaxed budget contraint}$$

Can I apply this to $N < \infty$?

$\sum_s a_n(t) \leq \alpha$		$\sum_s \mathbb{E}\left[a_n(t)\right] \leq \alpha$
Original problem (Hard)	$\longrightarrow$	LP relaxation (Easy)
V_N^*	$\leq$	V_{rel}^*

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Original problem
(Hard)

$$\longrightarrow$$

$$\sum_s \mathbb{E}\left[a_n(t)\right] \leq \alpha$$

LP relaxation
(Easy)

$$V_N^* \qquad \leq \qquad V_{rel}^*$$

$$\downarrow$$

$$x^*,\ y^*$$

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Original problem (Hard) $\longrightarrow$ LP relaxation (Easy)

$$V_N^* \qquad \leq \qquad V_{rel}^*$$

$$\wr\wr$$

$$V_N^\pi \quad \text{Can we build } \pi \text{ that is close to optimal?} \quad x^*,\ y^*$$

Main difficulty: in general $\mathbf{X}^N(t) \neq \mathbf{x}^*(t)$.

- We cannot choose $\mathbf{Y}^N(t) = \mathbf{y}^*(t)$.

Some historical perspective

- Infinite horizon: Index policies (Gittins 60s, Whittle index (98), Nino-Mora, 90s-2000s)
  - Often asymptotically optimal. (Weber and Weiss 91).

1. When they are: exponentially fast. (G, Gaujal, Yan 2021).
2. We can compute index efficiently. (G, Gaujal, Khun 2022).

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- Finite horizon: LP-index
  - Priority rule not always asymptotically optimal (Brown and Smith 2019), (Frazier et al 2020).

  3. When they are: exponentially fast (G, Gaujal, Yan 2022)

Outline

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Penalty and indexability

The $N = \infty$ is a constraint MDP:

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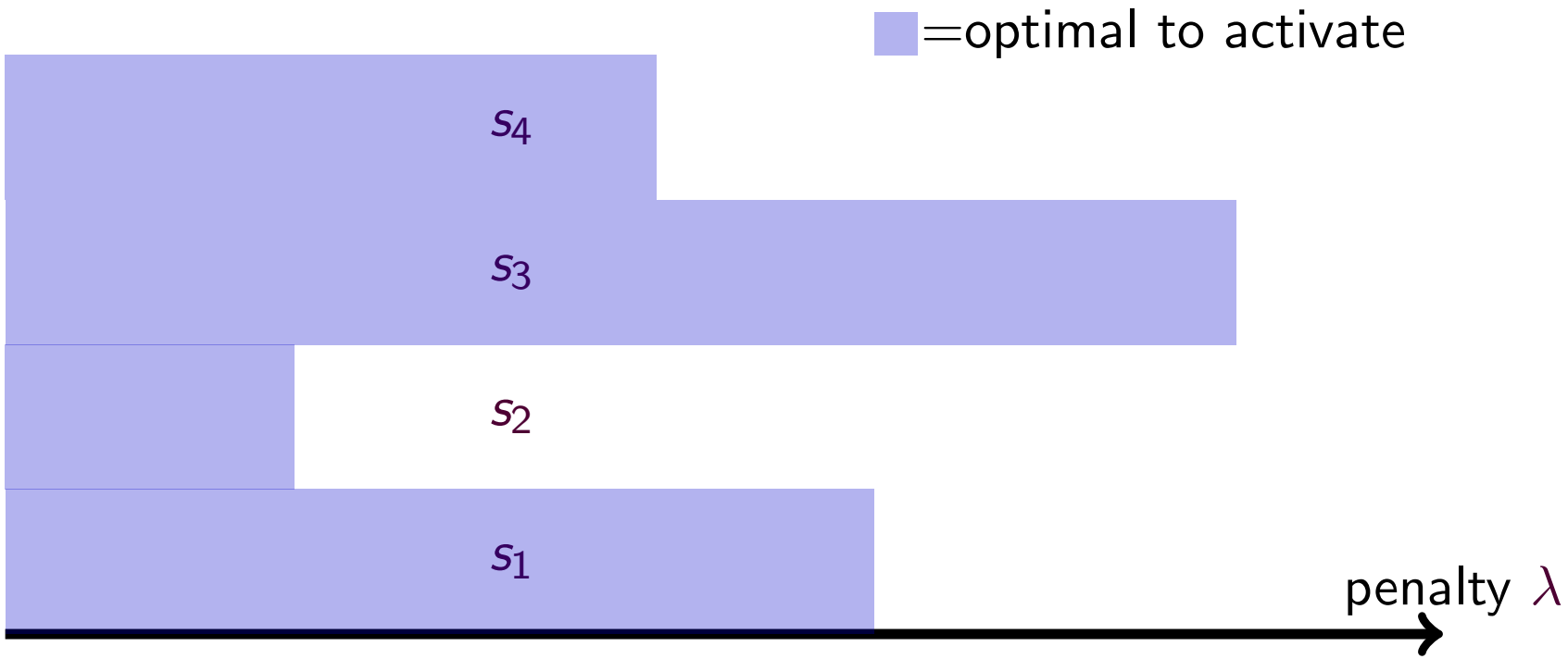
Idea: use a Lagrangian relaxation:

- $P(\cdot|s_n, a_n)$ and $r(s_n, a_n) - \lambda\, a_n$.

$\downarrow$

Penalty for activation

The penaly can be used to define a priority policy



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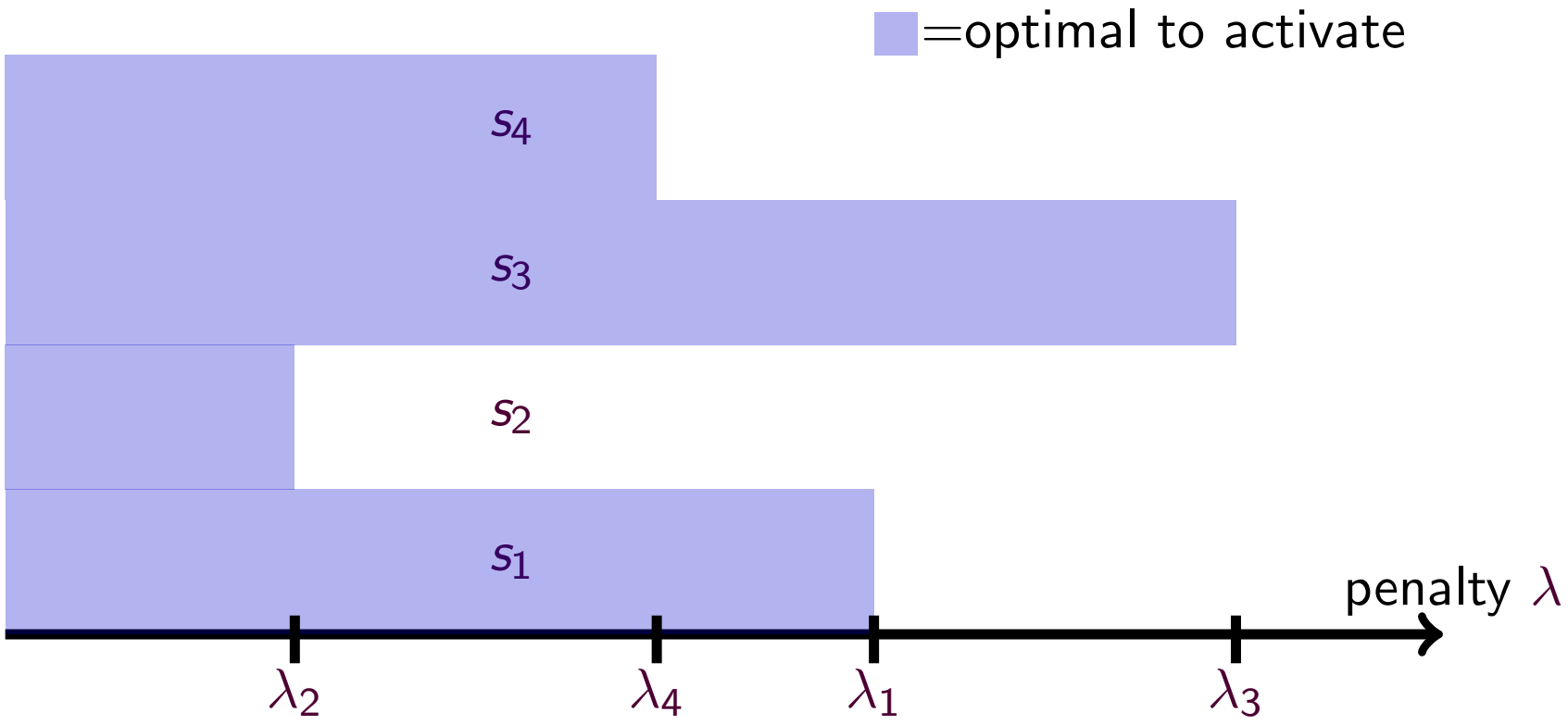

This is Whittle index policy.

For this example: $s_3 \succ s_1 \succ s_4 \succ s_2$.

Definition of Whittle index

Intuitively, for each state, there exists a λ_s such that any optimal policy is such that:

- The optimal action in s is 0 (rest) if $\lambda < \lambda_s$;
- The optimal action in s is 1 (activate) if $\lambda > \lambda_s$.

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This is not always true[1].

If the model satisfies this assumption, we say that the model is indexable. Whittle index policy is the corresponding priority policy.

[1] True with high probability? Yes: (Nino-Mora 01), No (G, Gaujal, Khun 21).

Illustration of what is Whittle policy

(stochastic scheduling)

Jobs of sizes X and Y with:

- $X = 10$
- $Y = \begin{cases} 2 & \text{proba } 1/2 \\ 18 & \text{proba } 1/2 \end{cases}$

Who should you run first to minimize expected completion time?

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Running a job costs 1€/sec and you can stop anytime. If you finish the job, you earn x. Whittle (=Gittins) index is the smallest x so that you start running the job.

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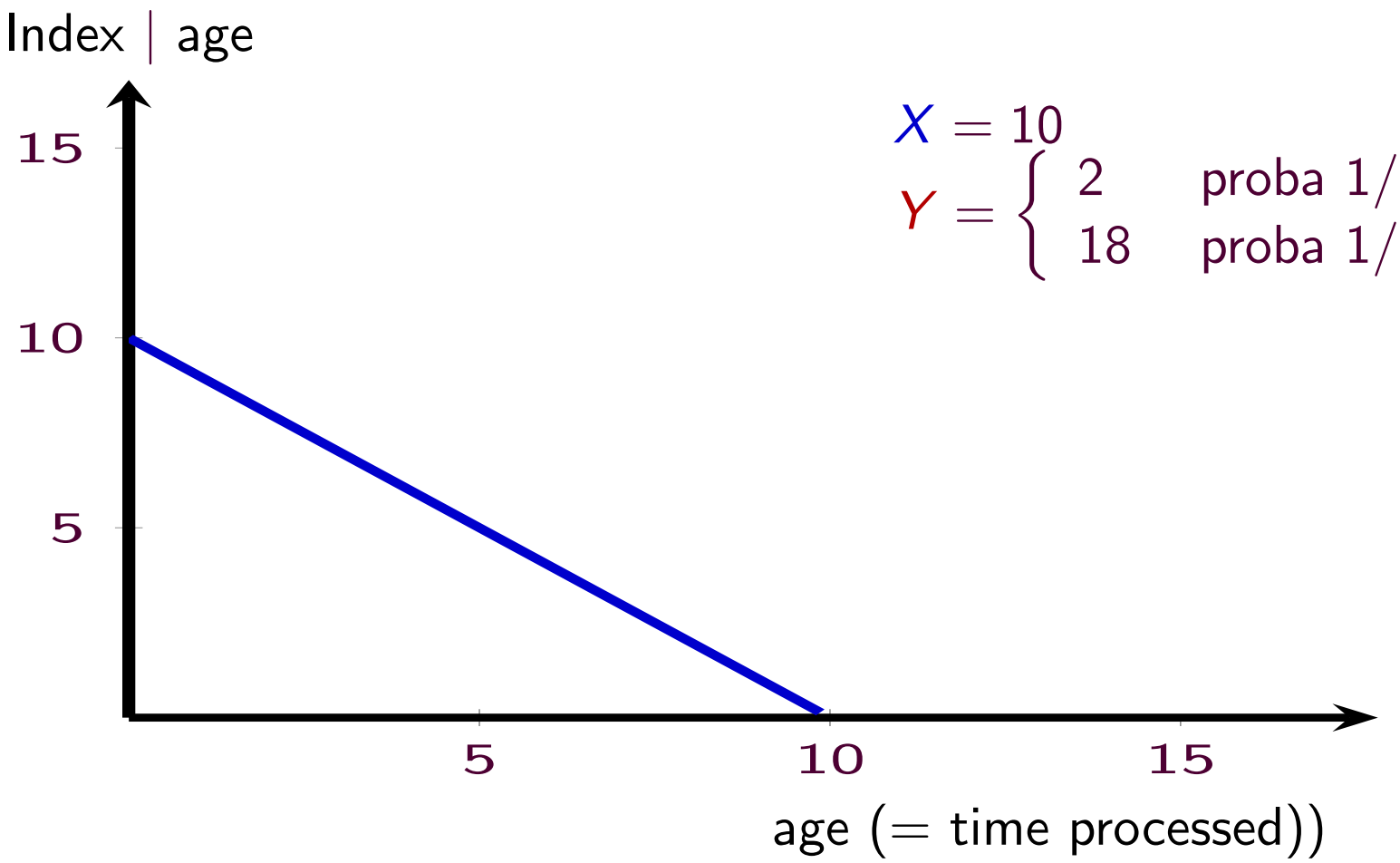

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Index can be computed independently for each job (=arm).

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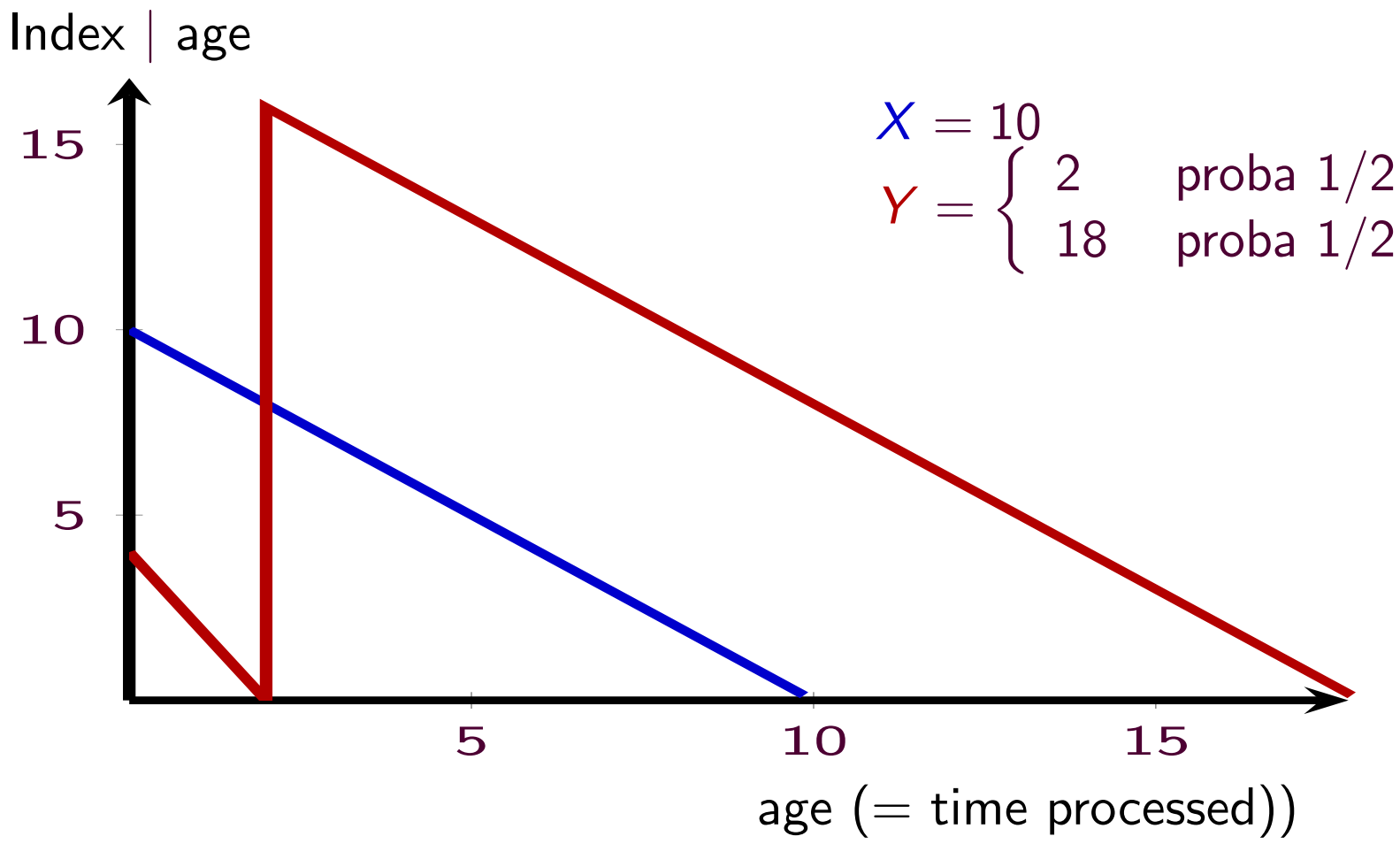

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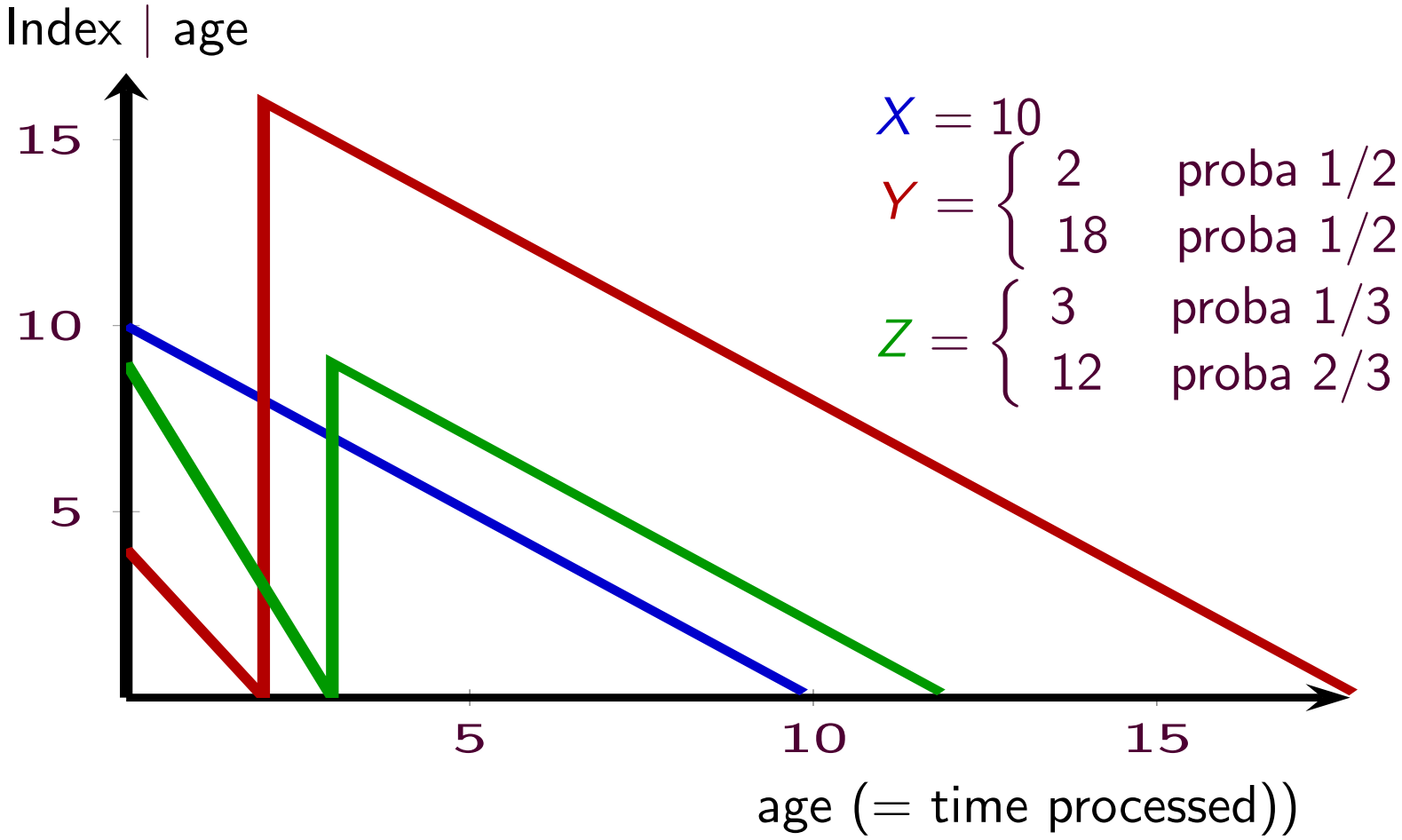

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Are Whittle index asymptotic optimal?

Assume indexability. For the infinite model, π^{WIP} defines a (piecewise linear) dynamical system:

$$\mathbf{x}(t+1) = \pi^{WIP}(\mathbf{x}(t)).$$

Theorem

1. *If π^{WIP} has a unique attractor, then WIP is asymptotically optimal.* *[Weber Weiss 90s, Verloop 2016]*
2. *For these problems, the suboptimality gap is exponentially small for non-degenerate problems.* *[G. Gaujal Yan 2021]*

Sketch of proof

Recall that $X_s^{(N)}(t) = \frac{1}{N}\#\{\text{arms in state } s \text{ at time } t\}$.

We have:

$$\mathbf{X}^{(N)}(t+1) = \pi^{WIP}(\mathbf{X}^{(N)}(t)) + \underbrace{O(1/\sqrt{N})}_{\text{stochastic noise. CLT}} .$$

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Hence:

1. If π^{WIP} has a unique attractor x^*, then $\mathbf{X}^N(\infty)$ concentrates on x^* (Hoeffding bound / large deviation).
2. Non-degenerate $=$ π^{WIP} is locally linear around x^*. We use the linearity of expectation.

How to compute Whittle indices?

Classical definition:

- The index is the penalty λ_i such that that an optimal policy can choose to activate or not the state i when the penalty is λ_i.

How to compute Whittle indices?

Refined definition:

- The index is the (unique) penalty λ_i such that that an (Bellman-)optimal policy can choose to activate or not the state i when the penalty is λ_i.

How to compute Whittle indices?

A Bellman-optimal policies satisfies Bellman equations:

$$g^*(\lambda) + h_i^*(\lambda) = \max_a \underbrace{r(i,a) + a\lambda + \sum_j P(j|i,a) h_j^*(\lambda)}_{q_{I,a}(\lambda)}$$

We define the *active advantage* $b_s(\lambda) := q_{i,1}(\lambda) - q_{i,0}(\lambda)$.

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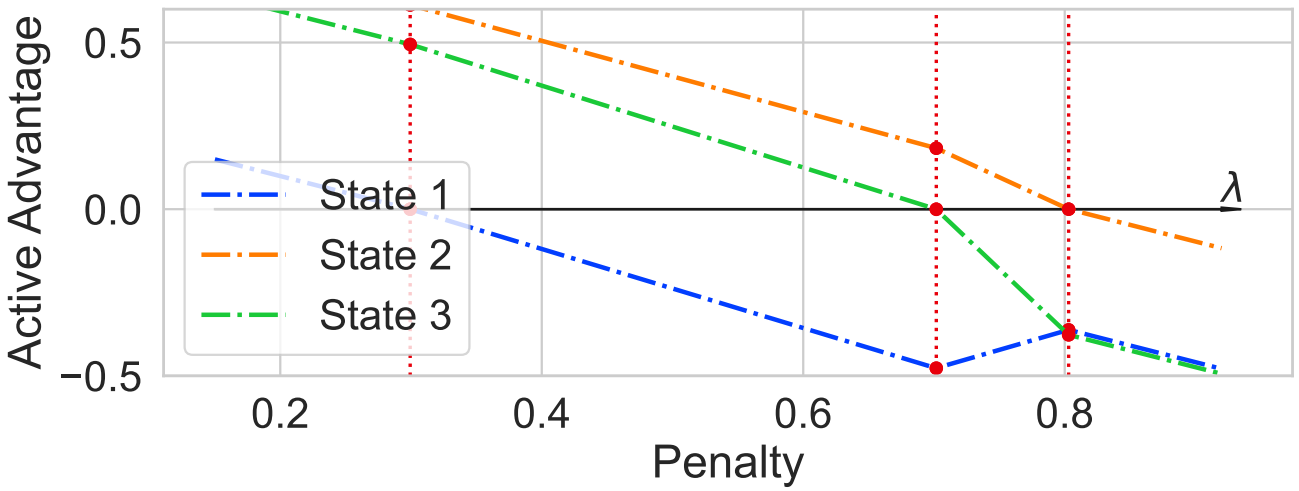
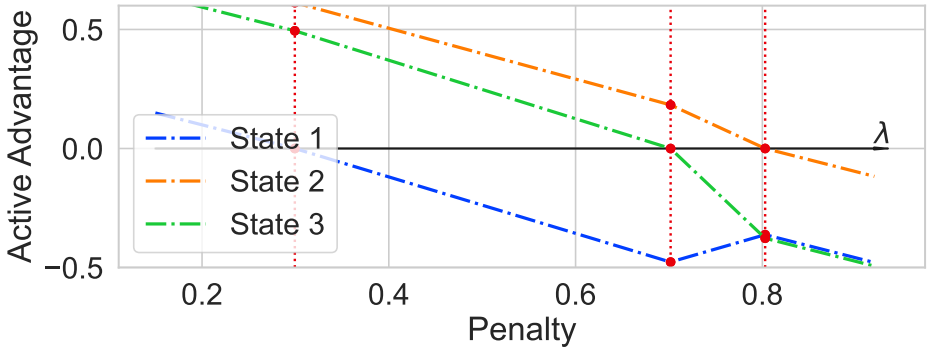


Theorem (G,Gaujal,Khun, 22)

An arm is indexable if and only if for all s: $b_{s,1}(\lambda) = 0$ has a unique solution.

We can use this characterization to build an efficient algorithm



Three ingredients:

1. For MDP, the advantage function is piecewise linear:

$$b^{\pi}(\lambda) = (A^{\pi})^{-1}(r + \lambda\pi).$$

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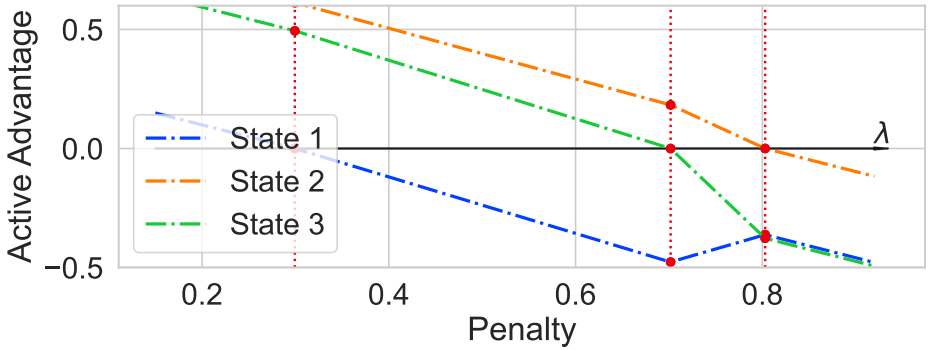

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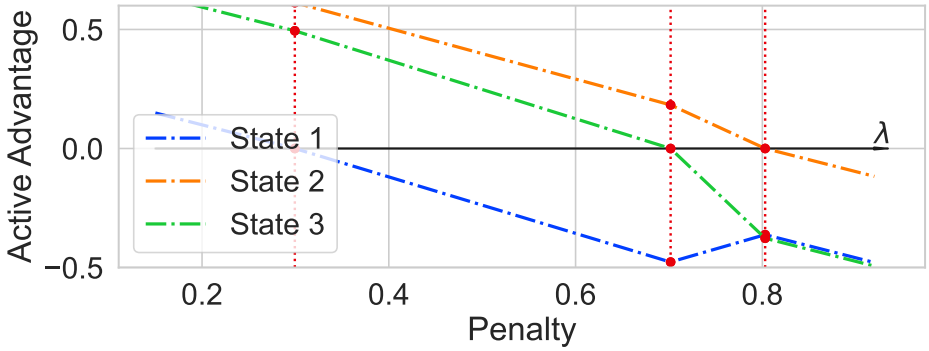

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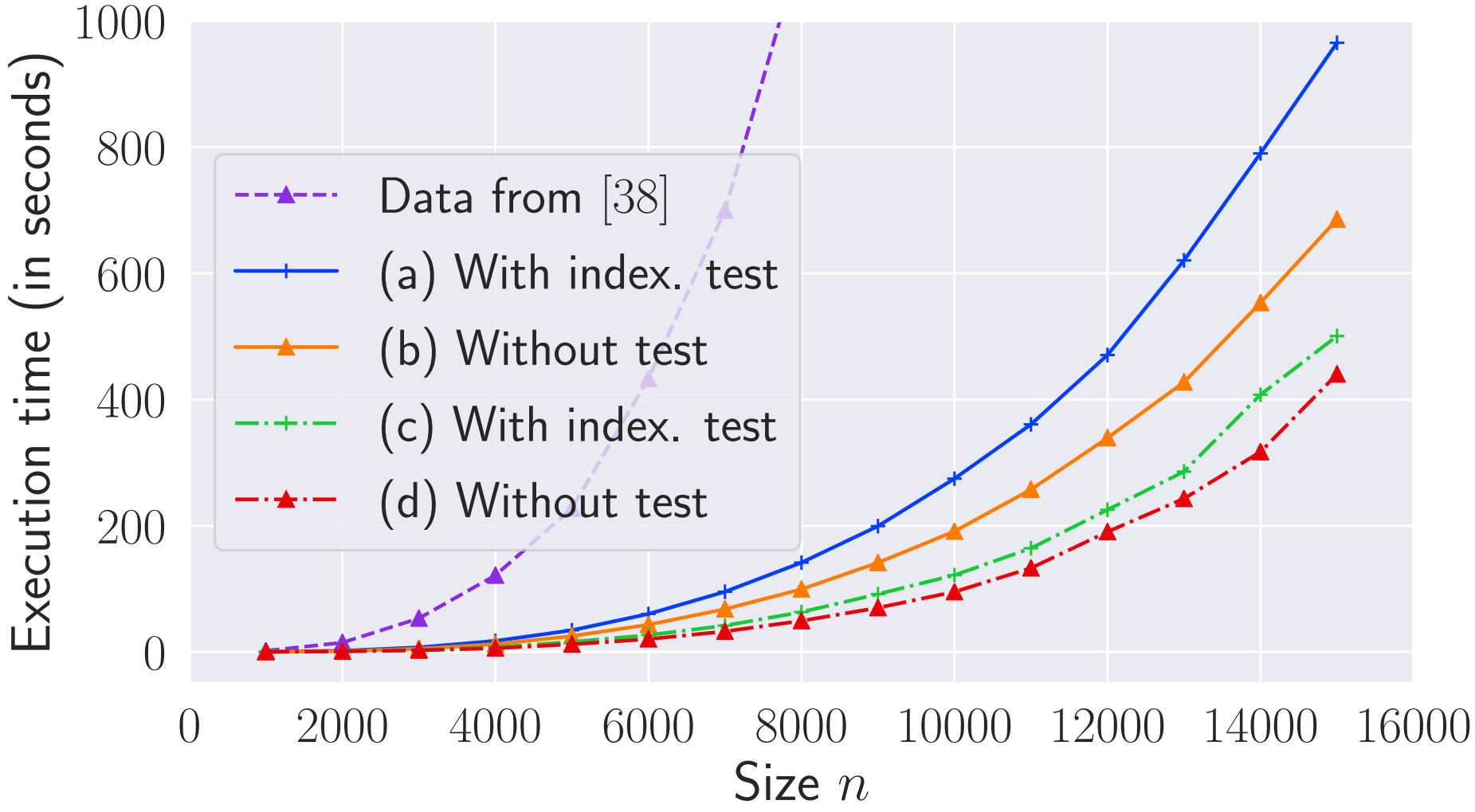
3. We can reorder operations to use Strassen's like operations.

We obtain a theoretical complexity of $O(S^{2.53})$ and an efficient implemenation

https://pypi.org/project/markovianbandit-pkg/



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How to construct a policy for the original problem?

Find a function $\pi : \mathcal{X} \to \mathcal{Y}$ such that $\pi(x^*) = y^*$.

Relaxed problem: Optimal sequence $x_s^*(t)$, $y_{s,a}^*(t)$.

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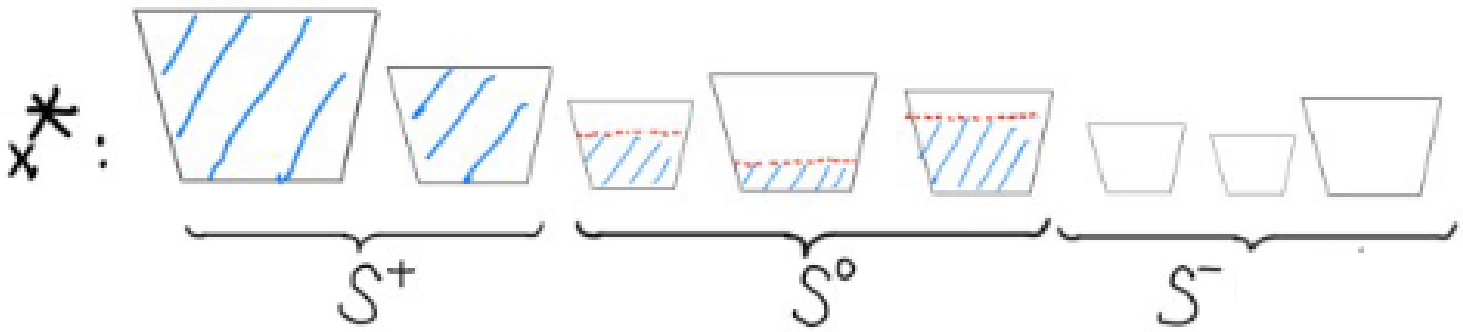

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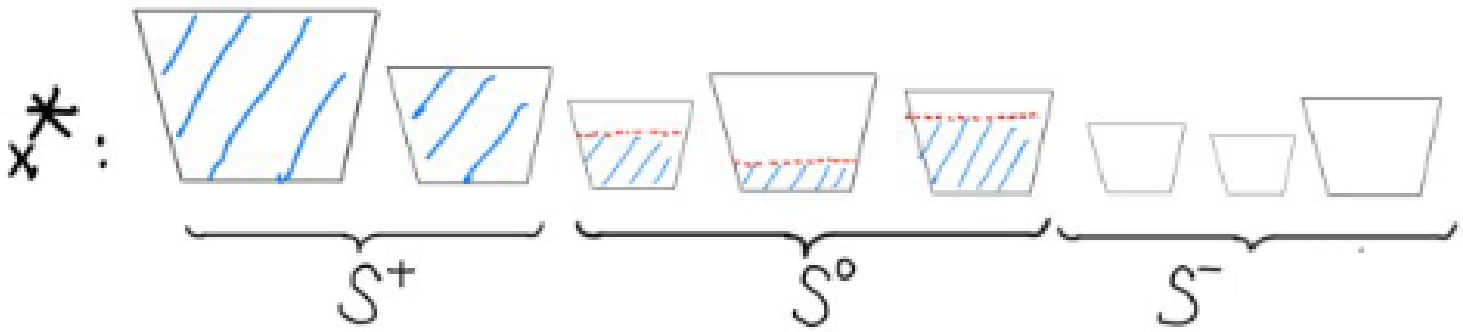
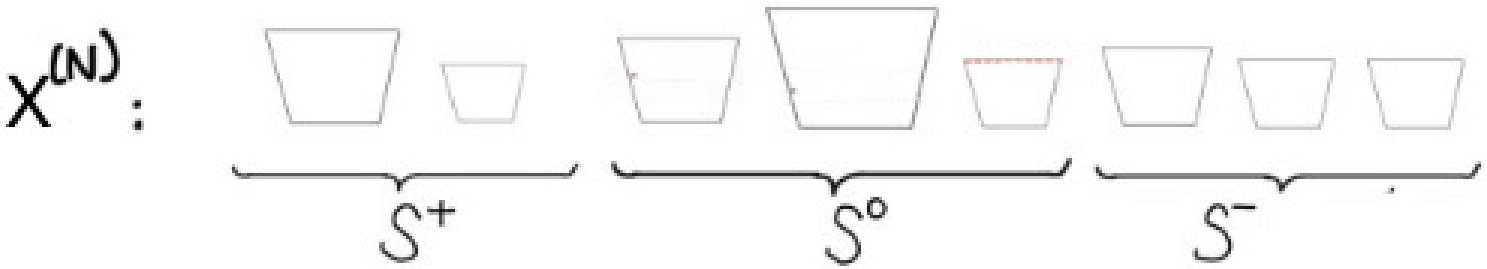


Original problem: Sequence $\pi_t : \mathcal{X} \to \mathcal{Y}$.



If $|calS^0(t)| = 1$, you can implement π_t as a priority rule.

- It is locally linear.

Asymptotic optimality

Theorem

- *There exists an priority rule that is asymptotically optimal if and only if for all t, $|S^0(t)| \leq 1$.*
- *It becomes optimal exponentially fast if for all t, $|S^0(t)| = 1$.*

Proof ingredients.

1. Concentration argument: π continuous implies $\lim_{N\to\infty} X_\pi^{(N)}(t) = x_\pi(t)$.
2. Linearity of expectation.

Many finite-horizon problems do not admit asymptotically optimal priority rules

Example: Applicant screening problem (Brown Smith 2020)

Candidates with prior quality Beta(1,1), Interview budget α=0.25

(1,1)

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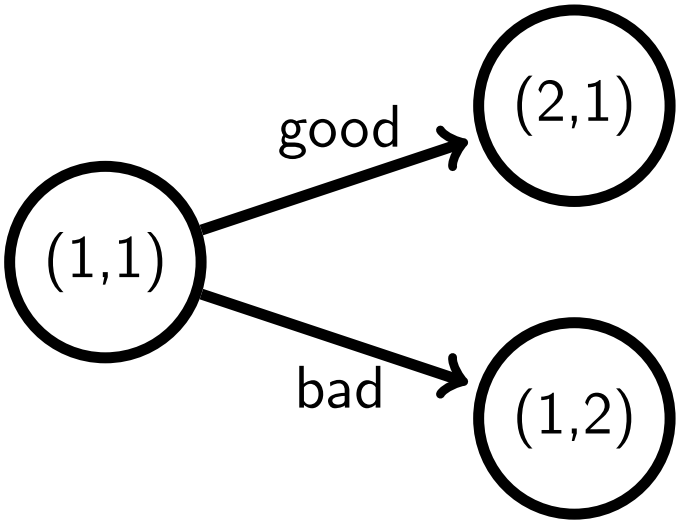

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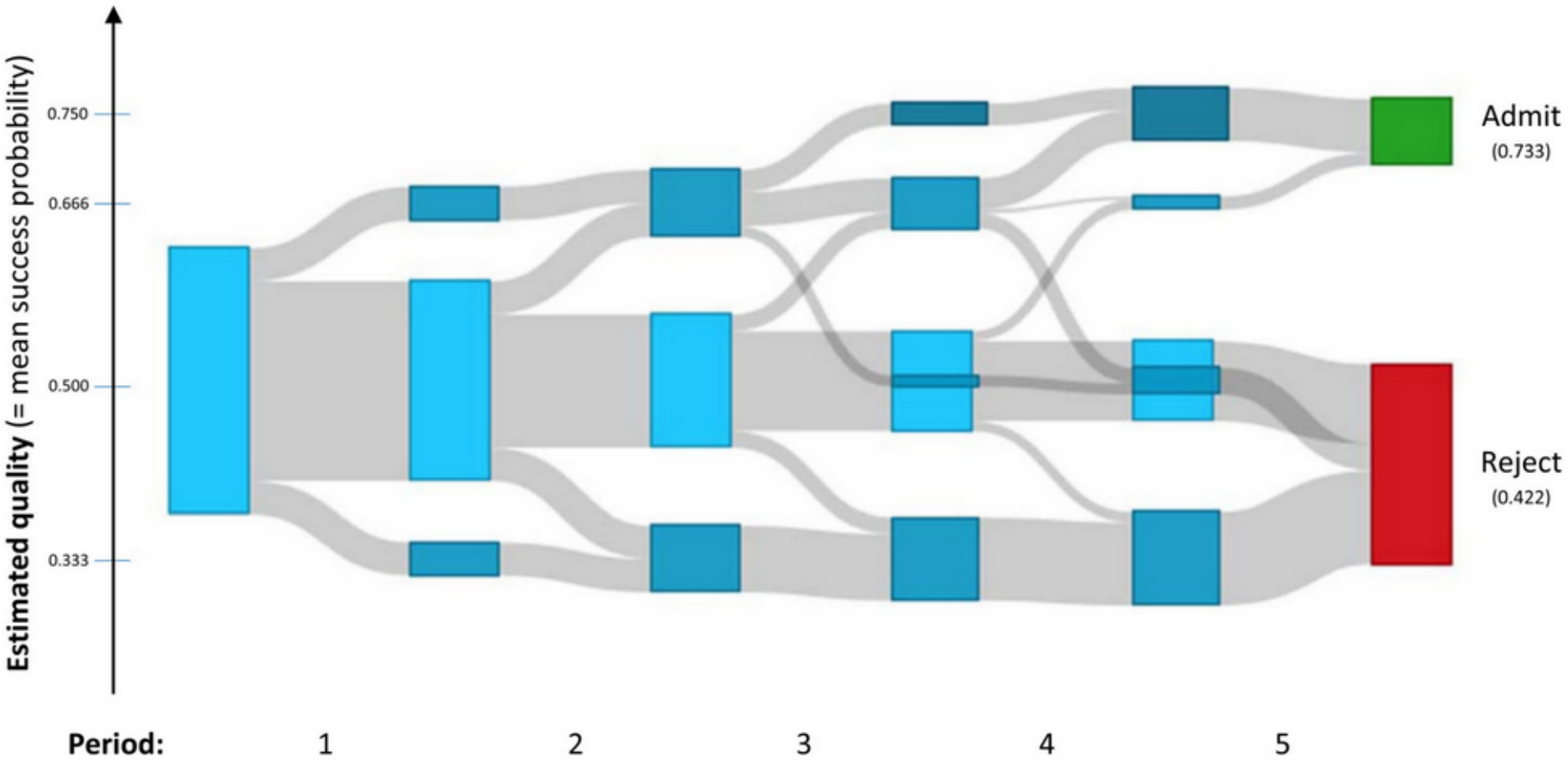

No asymptotically optimal priority policy: after two interviews:

- $x^*_{(1,1)} = 2x^*_{(2,1)} = 2x^*_{(1,2)} = 0.5$.
- $y^*_{(2,1),interview} = y^*_{(1,1),interview} = 0.125$.

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For Markovian bandits, mean-field control can be solved by an LP.

- Can be generalized to weakly coupled MDPs.

Simple policies (priority rule) are not always optimal.

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- Index policy (= "right actication price") are very efficient.

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- This talk: finite-state space, computation of policies.
- Open questions: learning, continuous state-spaces.

http://polaris.imag.fr/nicolas.gast/

- *Computing Whittle (and Gittins) Index in Subcubic Time*, G. Gaujal, Khun https://arxiv.org/abs/2203.05207
- *LP-based policies for restless bandits: necessary and sufficient conditions for (exponentially fast) asymptotic optimality*. G. Gaujal Yan. https://arxiv.org/abs/2106.10067