

Mean Field Approximation of Uncertain Stochastic Models

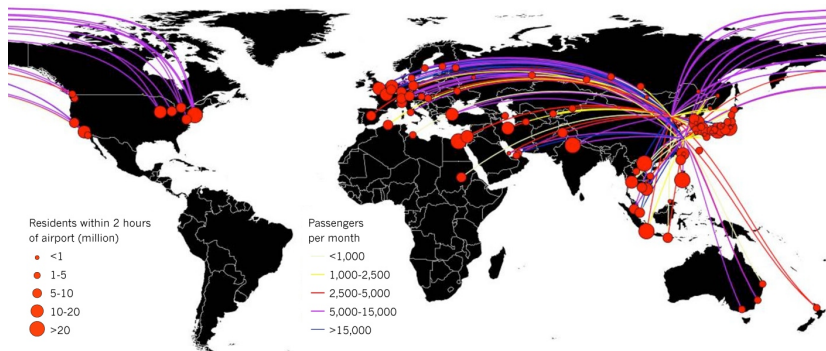
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DSN conference 2016, Toulouse, France

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Why do we need models?



- Design, prediction, optimization, correctness, etc.

Uncertainties in models: stochastic v.s. non-determinism

Stochastic
ex: Markov chains

Non-determinism
ex: ODE, timed-automata

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Quantitative analysis.

Worst case / correctness

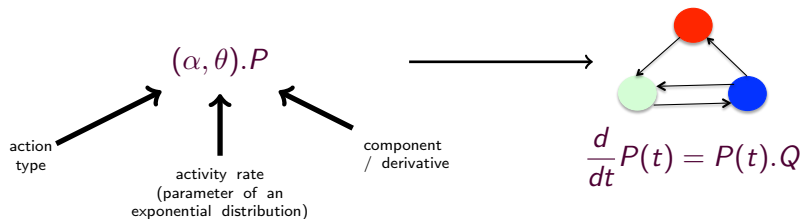
- + can be simulated
- How to choose parameters?

- Symbolic computation
- + No problem of parameters

Combination of the two

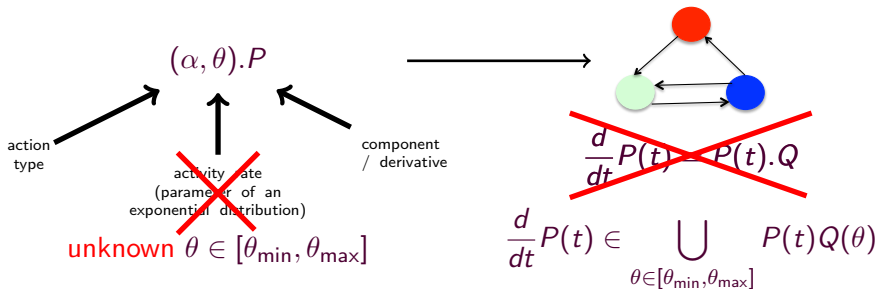
Uncertain Markov chains

- agents engage in actions at some rate.

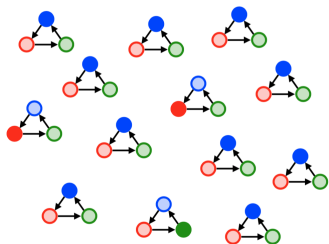


Building an uncertain continuous time Markov chain

- agents engage in actions at some rate.



Main problem: the state space grows exponentially



$3^{13} \approx 10^6$ states.

We need to keep track

$$\mathbb{P}(X_1(t) = i_1, \dots, X_n(t) = i_n)$$

and solve the differential inclusion:

$$\frac{d}{dt}P(t) \in \bigcup_{\theta \in [\theta_{\min}, \theta_{\max}]} P(t)Q(\theta)$$

Is there any hope?

Contributions (and Outline)

- ① Some systems simplify when the population grows.
 - ▶ Mean-field approach
- ② We can add non-determinism to these models
- ③ We can build and use numerical algorithms.

Outline

- 1 Population Processes and Classical Mean Field Methods
- 2 Uncertain and Imprecise Population Processes
- 3 Numerical Algorithms and Comparisons
 - Numerical algorithms (transient regime)
 - Steady-state
 - General processor sharing example
- 4 Conclusion

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Mean field methods have been used in a multiple contexts

ex: model-checking, performance of SSD, load balancing, MAC protocol,...

SPAA 98 **Analyses of Load Stealing Models Based on Differential Equations** by Mitzenmacher

JSAC 2000 **Performance Analysis of the IEEE 802.11 Distributed Coordination Function** by Bianchi

FOCS 2002 **Load balancing with memory** by Mitzenmacher et al.

DSN 2013 **A logic for model-checking mean-field models** by Kolesnichenko et al

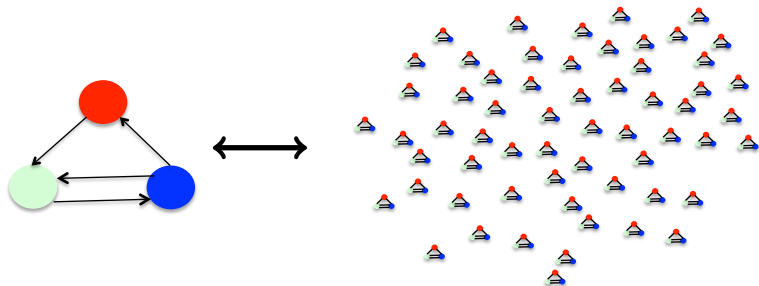
DSN 2013 **Lumpability of fluid models with heterogeneous agent types** by Iacobelli and Tribastone

SIGMETRICS 2013 **A mean field model for a class of garbage collection algorithms in flash-based solid state drives** by Van Houdt

⋮ ⋮

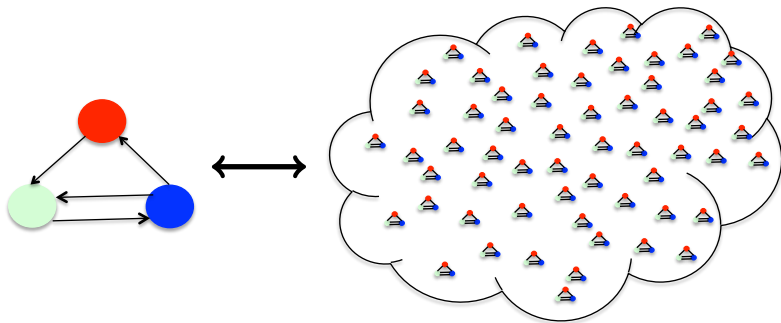
These models correspond to **distributed** systems

Each object interacts with the mass



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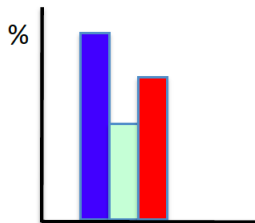
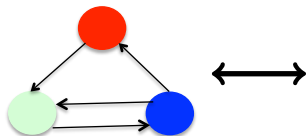
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We view the population of objects more abstractly, assuming that individuals are indistinguishable.

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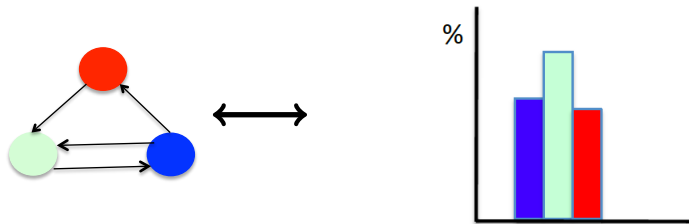


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An **occupancy measure** records the proportion of agents that are currently exhibiting each possible state.

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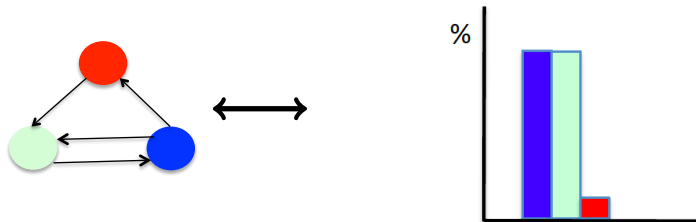


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Population CTMC

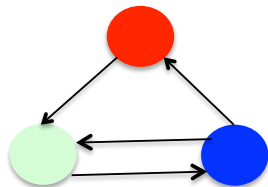
We consider a sequence \mathbf{X}^N , indexed by the population size N , with state spaces $\mathbf{E}^N \subset E \subset \mathbb{R}^d$. The transitions are:

$$X \mapsto X + \frac{\ell}{N} \quad \text{at rate } N\beta_\ell(X).$$

for a finite number of $\ell \in \mathcal{L}$.

The drift is $f(x) = \sum_{\ell} \ell \beta_\ell(x)$.

Example :



The state is (X_S, X_I, X_R) and the transitions are

$$\ell_1 = (-1, +1, 0) \text{ at rate } NX_S X_I$$

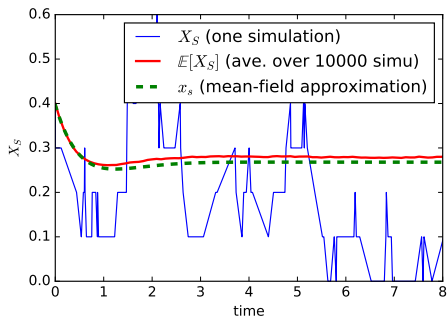
$$\ell_2 = (0, -1, +1) \text{ at rate } NX_I$$

$$\ell_3 = (+1, 0, -1) \text{ at rate } NX_R$$

$$\ell_4 = (-1, 0, +1) \text{ at rate } NX_S$$

Kurtz' convergence theorem

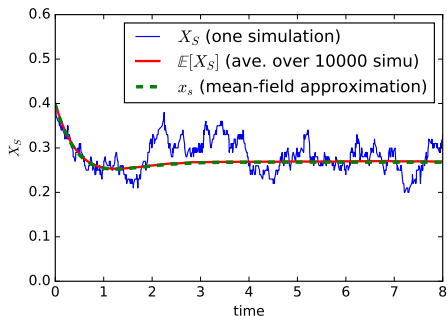
Theorem: Let \mathbf{X} be a population model. If $X^N(0)$ converges (in probability) to a point x , then the stochastic process \mathbf{X}^N converges (in probability) to the solutions of the differential equation $\dot{x} = f(x)$, where f is the drift.



$$N = 10$$

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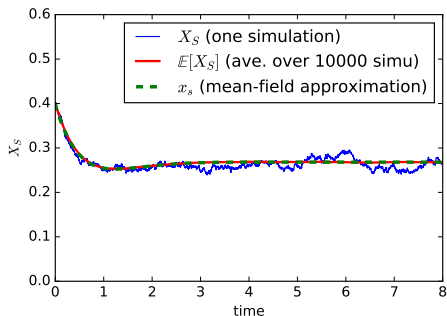
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Uncertain and imprecise models

Instead of $\beta_\ell(x)$, the rates can depend on a parameter ϑ : $\beta_\ell(x, \vartheta)$.

We distinguish two kinds of uncertainties:

Uncertain

$\vartheta \in \Theta$ is constant but its value is not known precisely.

- Uncertainties in the model

Imprecise

$\vartheta = \vartheta(t) \in \Theta$ can vary (measurably) as a function of time

- human behavior, environment, adversary,...

Uncertain and imprecise population models

We consider a sequence \mathbf{X}^N of Imprecise or Uncertain population processes, indexed by the size N , with state spaces $\mathbf{E}^N \subset E \subset \mathbb{R}^d$. The transitions are (for $\ell \in \mathcal{L}$):

$$X \mapsto X + \frac{\ell}{N} \quad \text{at rate } N\beta_\ell(X, \vartheta)$$

The drifts corresponding to parameter ϑ is $f(x, \vartheta) = \sum_{\ell \in \mathcal{L}} \ell \beta_\ell(x, \vartheta)$.

Theorem (Bortolussi, G. 2016)

if $X^N(0)$ converges (in probability) to a point x , then the uncertain (or imprecise) stochastic process \mathbf{X}^N converges in probability to:

<i>Uncertain</i>	<i>Imprecise</i>
<i>A solution of $\dot{x} = f(x, \vartheta)$ (for a given ϑ)</i>	<i>A solution of $\dot{x} \in \bigcup_{\vartheta \in \Theta} f(x, \vartheta)$</i>

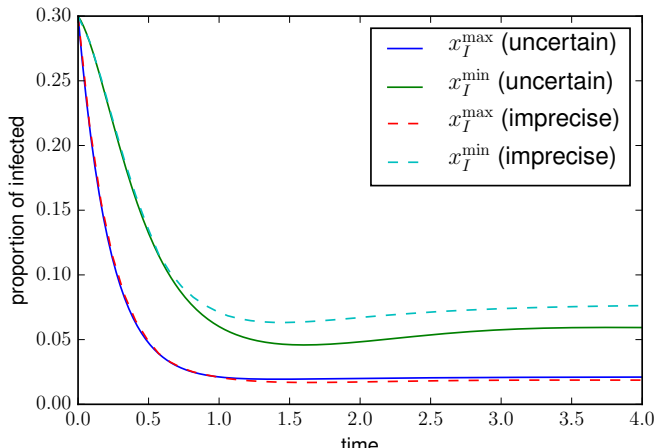
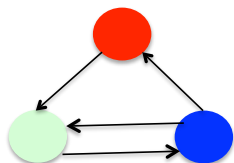
Example of the SIR model

The state is (X_S, X_I, X_R) and the transitions are

$$l_1 = (-1, +1, 0) \text{ at rate } N(aX_S + \vartheta X_S X_I)$$

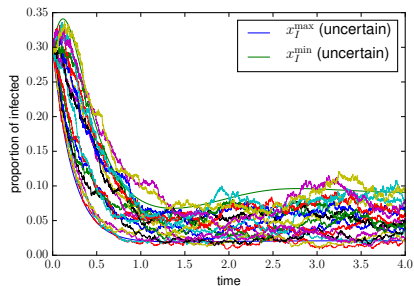
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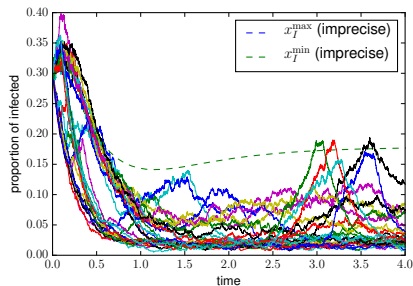


Consequence on the stochastic system

Uncertain



Imprecise

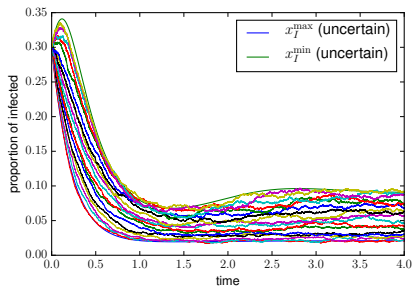


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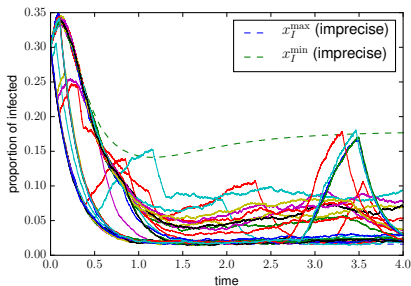
Remark : the proportion of infected is non-monotone on the infection rate.

Consequence on the stochastic system

Uncertain



Imprecise



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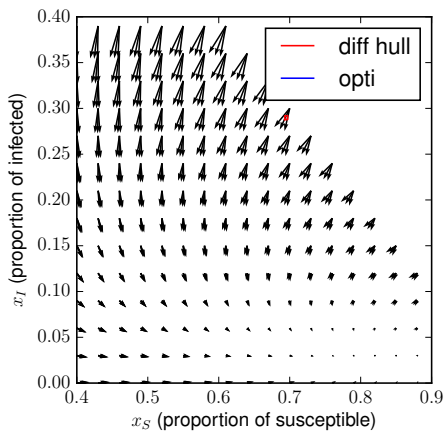
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Numerical algorithms

- Uncertain (fixed parameter)
 - ▶ Exhaustive search
 - ▶ Online learning
- Imprecise (varying parameter). Difficulty = non-linear.
 - ▶ Exact: reachability (ex: solvable by Pontryagin's principle)
 - ▶ Approximation: polygons (ex: differential hull)

Example of numerical algorithm: differential hull



$$\vartheta \in \Theta = [1, 3]$$

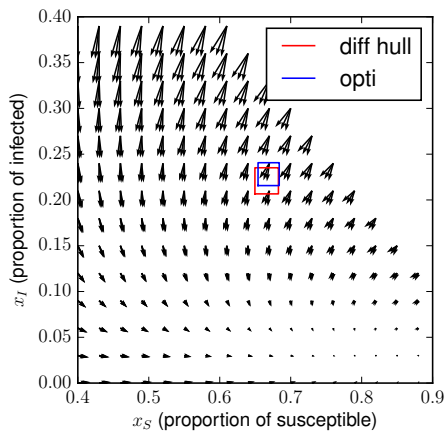
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$$\underline{x}(t) \leq x(t) \leq \bar{x}(t)$$

where \underline{x} and \bar{x} satisfy $\dot{\underline{x}} = \underline{f}(\underline{x}, \bar{x})$ and $\dot{\bar{x}} = \bar{f}(\underline{x}, \bar{x})$, with

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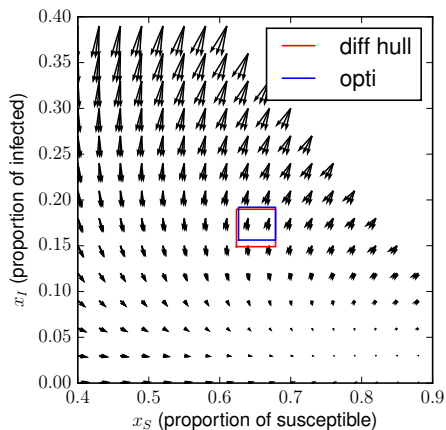
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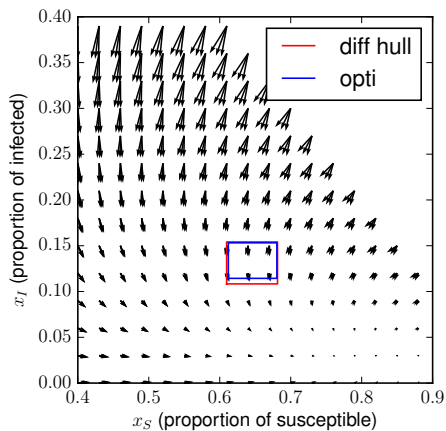
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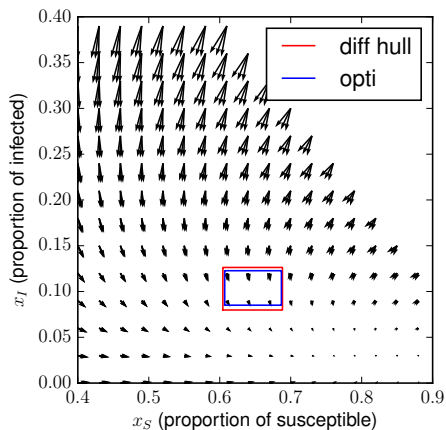
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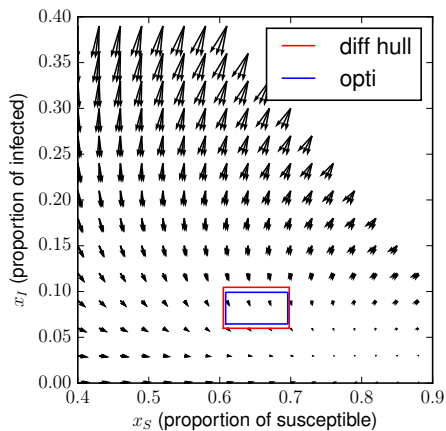
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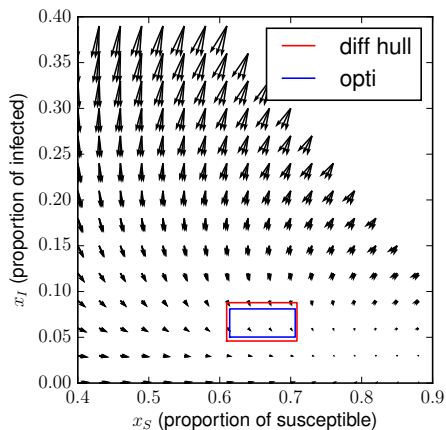
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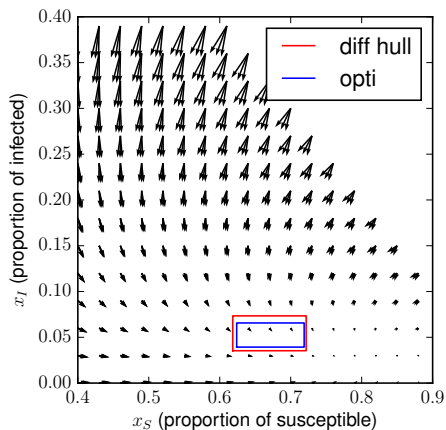
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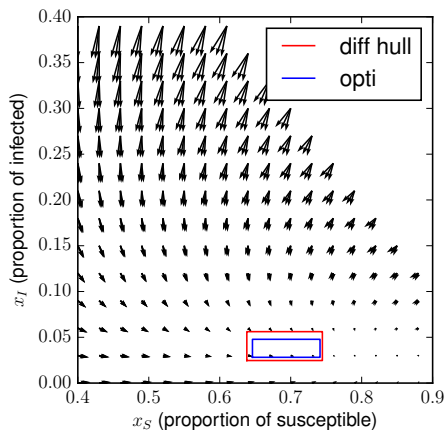
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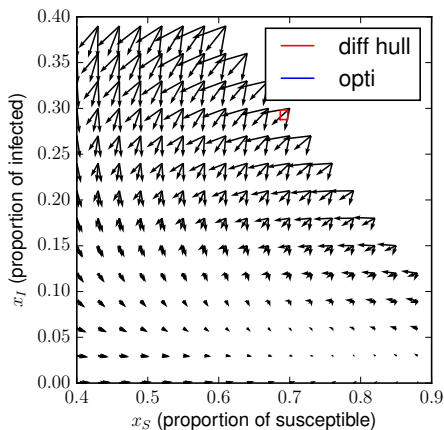
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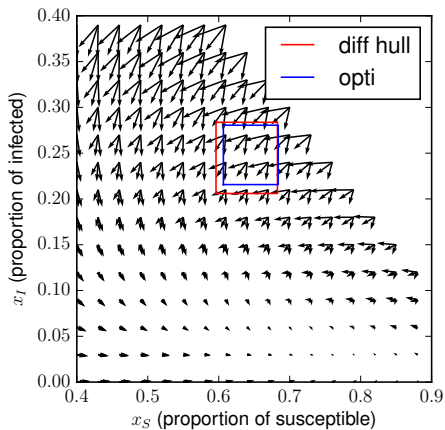
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Differential hull provide loose bounds when Θ is large.

Example of numerical algorithm: differential hull



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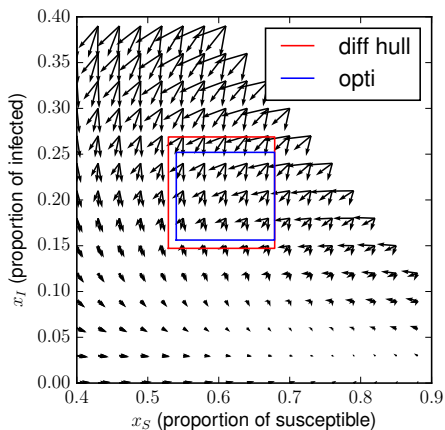
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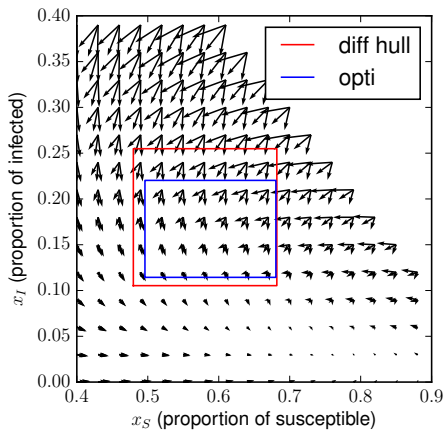
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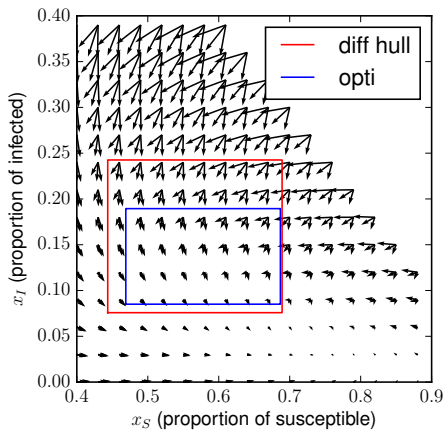
where \underline{x} and \bar{x} satisfy $\dot{\underline{x}} = \underline{f}(\underline{x}, \bar{x})$ and $\dot{\bar{x}} = \bar{f}(\underline{x}, \bar{x})$, with

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Differential hull provide loose bounds when Θ is large.

Example of numerical algorithm: differential hull



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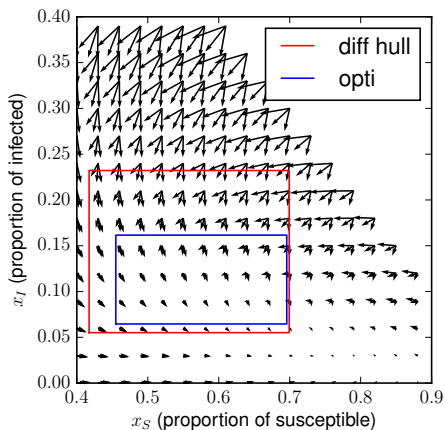
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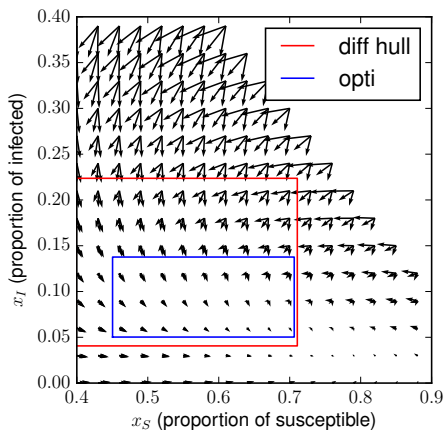
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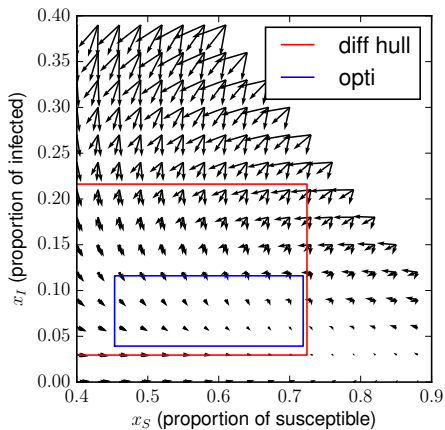
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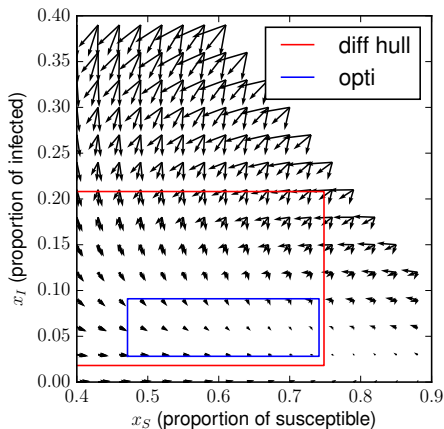
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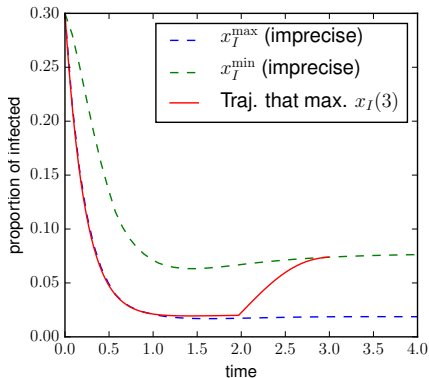
Differential hull provide loose bounds when Θ is large.

An alternative is to formulate the problem as an optimization problem

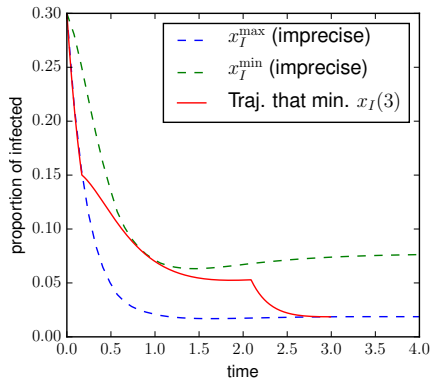
$$x_i^{\max}(T) := \max_{\theta} x_i(T) \text{ such that for all } t \in [0; T]:$$
$$\begin{cases} x(t) = x + \int_0^t f(x(s), \theta(s)) ds \\ \theta(t) \in \Theta \end{cases}$$

- Pontryagin's principle can be used.
- Easier than MDP

SIR model: we can compute an minimal/maximal trajectory by using Pontryagin's maximum principle



A trajectory that maximizes $x_I(3)$.



A trajectory that minimizes $x_I(3)$.

Stationary regime of imprecise models

Stationary regime of imprecise models

- Asymptotic reachable set of the differential inclusion A_F :

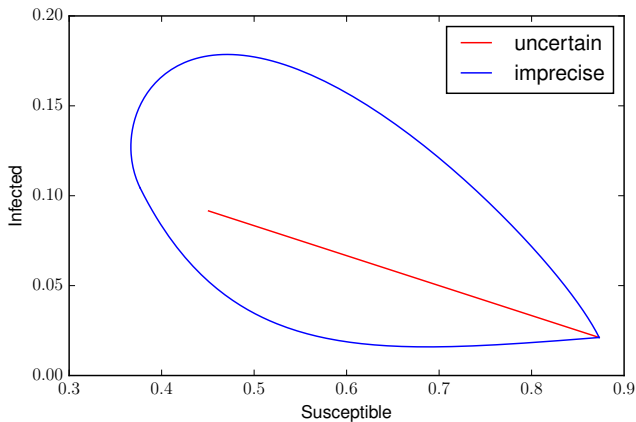
$$A_F = \bigcap_{T>0} \bigcup_{x, t \geq T, x \in S_{F,x}} \{\mathbf{x}(t)\}$$

- Theorem:** Let \mathbf{X} be an imprecise population process, then

$$\lim_{N \rightarrow \infty} \lim_{t \rightarrow \infty} d(X^N(t), A_F) = 0 \quad \text{in probability.}$$

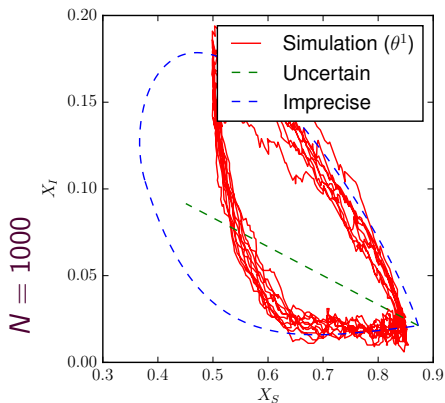
- Theorem:** Let \mathbf{X} be an imprecise population process such that \mathbf{X}^N is a Markov chain that has a stationary measure μ^N . Let μ be a limit point of μ^N (for the weak convergence). Then, the support of μ is included in the **Birkhoff centre** of F : $\mu(B_F) = 1$.

Asymptotically reachable set: example for the SIR model

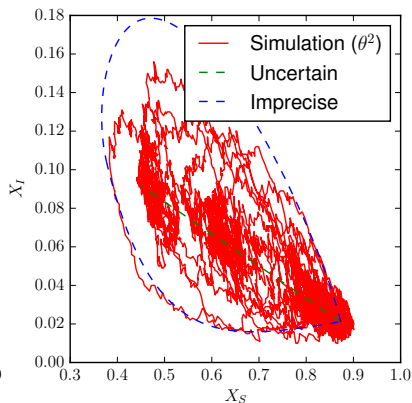


Note: comparison with the differential hull approach: bounds are very loose.

SIR model: stationary regime



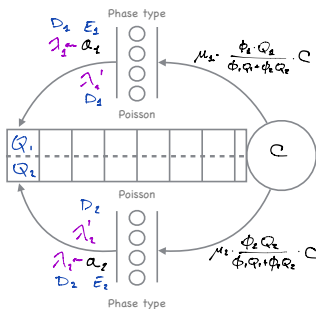
(a) policy θ^1



(b) policy θ^2

- No policy can make the stochastic system exit the blue zone (for large N).

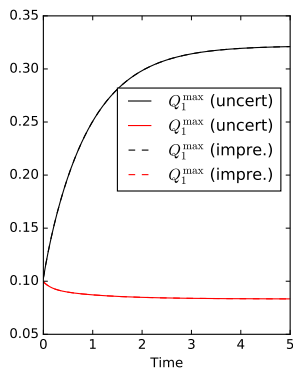
In the paper, we also study a Generalized Processor Sharing model



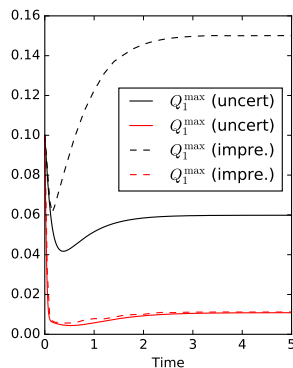
- Parameters: $\mu_1 = 5$, $\mu_2 = 1$, $\phi_1 = \phi_2 = 1$, $a_1 = 1$ and $a_2 = 2$. λ_i, λ_i' imprecise with $\lambda_1^{\min} = 1$, $\lambda_1^{\max} = 7$, $\lambda_2^{\min} = 2$, $\lambda_2^{\max} = 3$, $\lambda_i'^{\min} = 1/(1/a_i + a/\lambda_i^{\min})$, and $\lambda_i'^{\max} = 1/(1/a_i + a/\lambda_i^{\max})$

- A model of two tandem queues Q_1, Q_2 sharing a processor. Q_i gets a fraction $\phi_i N_i Q_i / (\phi_1 N_1 Q_1 + \phi_2 N_2 Q_2)$ of the capacity C of the server. Each queue serves a job of type i , with average completion time μ_i . Arrivals are Poisson (D_i - delay station - rate λ_i') or MAP (two delay stations in series E_i, D_i rates a_i, λ_i).

Generalized Processor Sharing: for the imprecise model, a higher arrival rate does not imply a larger queuing delay.



Poisson



MAP

- Optimization (for imprecise) of ϕ_1 to minimize the maximum queue length at time t : $\bar{Q}(t) = \max_{\theta} (Q_1^{\theta}(t) + Q_2^{\theta}(t))$.

Outline

- 1 Population Processes and Classical Mean Field Methods
- 2 Uncertain and Imprecise Population Processes
- 3 Numerical Algorithms and Comparisons
 - Numerical algorithms (transient regime)
 - Steady-state
 - General processor sharing example
- 4 Conclusion

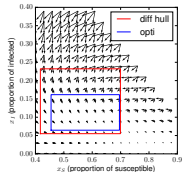
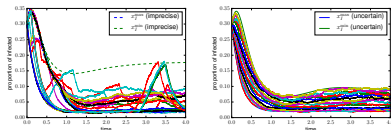
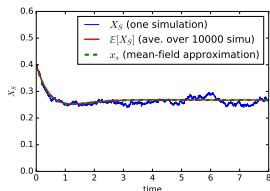
Recap and Future Work

Mean field methods are useful to study large stochastic systems.

We extended the mean field results for imprecise and uncertain PCTMCs, both at transient and at steady state.

We developed numerical method to bound the reachable sets.

Future work: scalability of numerical algorithms, integration in a toolset (EU project Quanticol), algorithmic complexity.



Thank you!

Slides are online:

<http://mescal.imag.fr/membres/nicolas.gast>

nicolas.gast@inria.fr

This paper *Mean Field Approximation of Uncertain Stochastic Models*, L. Bortolussi and N. Gast, DSN 2016

Other mean-field references

B-G 16 *Mean-Field Limits Beyond Ordinary Differential Equations*, L. Bortolussi, N. Gast., SFM Quanticol summer school

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G. Van Houdt 15 *Transient and Steady-state Regime of a Family of List-based Cache Replacement Algorithms.*, Gast, Van Houdt., ACM Sigmetrics 2015